CLAIMS:

- 1. A method for cryptographically converting an input data block into an output data block; the method including performing a non-linear operation on the input data block using an S-box based on a permutation, wherein the method includes each time before using the S-box (pseudo-)randomly selecting the permutation from a predetermined set of at least two permutations associated with the S-box.
- 2. A method as claimed in claim 1, wherein the set of permutations is formed such that a cryptographic weakness in one of the permutations of the set is at least partially compensated by a corresponding cryptographic strength in at least one of the other permutations of the set.
- 3. A method as claimed in claim 1, wherein the data block consists of n data bits and each element of the set of permutations is a permutation on a set of  $2^n$  elements, represented by  $\mathbb{Z}_2^n$ , where each non-trivial differential characteristic of each permutation in this set has a probability of at most  $p_{diff}$ , the set of permutations being formed by permutations which have been selected such that for each non-trivial differential characteristic with probability of  $p_{diff}$  in any of the permutations, this differential characteristic has a probability lower than  $p_{diff}$  in at least one of the other permutations of the set.

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- 4. A method as claimed in claim 3, wherein the differential characteristic has a probability equal to zero in at least one of the permutations.
- 5. A method as claimed in claim 4, wherein n = 4, and  $p_{diff} = \frac{1}{4}$ .

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6. A method as claimed in claim 1, wherein the data block consists of n data bits and each element of the set of permutations is a permutation on a set of  $2^n$  elements, represented by  $\mathbb{Z}_2^n$ , where each non-trivial linear characteristic of each permutation in this set has a probability of at least  $\frac{1}{2} - p_{lin}$  and at most  $\frac{1}{2} + p_{lin}$ , the set of permutations being formed

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by permutations which have been selected such that for each non-trivial linear characteristic with probability of  $\frac{1}{2}$  -  $p_{lin}$  or  $\frac{1}{2}$  +  $p_{lin}$  in any of the permutations, this linear characteristic has a probability closer to  $\frac{1}{2}$  in at least one of the other permutations of the set.

- 5 7. A method as claimed in claim 5, wherein the linear characteristic has a probability equal to ½ in at least one of the permutations.
  - 8. A method as claimed in claim 6, wherein n = 4 and  $p_{lin} = \frac{1}{4}$ .
- 10 9. A method as claimed in claim 1, wherein the set of permutations consists of two permutations.
  - 10. A method as claimed in claim 1, including performing the selection of the permutation under control of an encryption key.

11. A method as claimed in claim 9 and 10, wherein the selection of the permutation is performed under control of one bit of the encryption key.

- 12. A computer program product where the program product is operative to cause a processor to perform the method of claim 1.
  - 13. A system for cryptographically converting an input data block into an output data block; the method system including:
    - an input for receiving the input data block;
  - a storage for storing a predetermined set of at least two permutations associated with an S-box;
    - a cryptographic processor for performing a non-linear operation on the input data block using an S-box based on a permutation; the processor being operative to, each time before using the S-box, (pseudo-)randomly selecting the permutation from the stored set of permutations associated with the S-box; and
      - an output for outputting the processed input data block.